What is Graph- and Multi-Graph Matching and Why we need them

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What is Graph-Matching?





What is Graph-Matching?



Graph matching, linear/quadratic assignment problem, weighted bipartite matching

Applications of (Multi-)Graph-Matching in CV



Keypoint matching of different objects

Image courtesy: [Rolinek et al. 2020]



Non-rigid motion estimation

Image courtesy : [Alhaija et al. 2015]



Cell matching and tracking

Related work

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> 30 exact methods, > 100 heuristic algorithms

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2018: Zanfir, Sminchisescu. Deep learning of graph matching
2022: Haller et al. A comparative study of graph matching algorithms in computer vision

What we are going to address in the tutorial

13:30-14:00 What is graph- and multi-graph matching and why we need it

14:00-14:30 **Applications** in computer vision and bio-imaging

14:30-15:00 **Optimization** for graph matching and a recent comparison study

15:00 - 15:30 - Coffee break

15:30-16:15 Deep graph matching: **Learning** to match

ulti-graph matching: Optimization methods and applications











Outline (of this particular talk)

Linear and quadratic assignment problems

- Expressing power
- Integer programs
- Complexity

Where the name graph matching comes from?

- Classical formulation
- Koopmans-Beckmann form

How different is it in Computer Vision vs. Operations Research?

- Outliers and (in)complete matching
- Speed vs. precision

Multigraph matching

Definition and complexity









Another name: Weighted bipartite matching











$$\min_{x \in \{0,1\}^{n \times n}} \sum_{i=1}^{n} \sum_{s=1}^{n} c_{is} x_{is}$$

s.t.:
$$\sum_{s=1}^{n} x_{is} = 1 \ \forall i$$
$$\sum_{i=1}^{n} x_{is} = 1 \ \forall s$$

$$= \min_{\substack{x \ge 0}} \sum_{i=1}^{n} \sum_{s=1}^{n} c_{is} x_{is}$$

s.t.:
$$\sum_{s=1}^{n} x_{is} = 1 \quad \forall i$$
$$\sum_{i=1}^{n} x_{is} = 1 \quad \forall s$$





 c_{is} - cost of matching $i \leftrightarrow s$ $x_{is} \in \{0, 1\}$ - edge selectors $\sum_{is} c_{is} x_{is}$ - total matching cost

Efficiently solvable e.g. by Hungarian method O(n³)

Python: scipy.optimize.linear_sum_assignment()





appearance cost only, no geometry!



appearance cost only, no geometry!



Very similar appearance, geometry is crucial!





$$c_{is,jl} := (d_{ij} - d_{sl})^2$$



Total matching cost:

$$\sum_{is} c_{is} x_{is} + \sum_{\substack{i \neq j \\ s \neq l}} c_{is,jl} x_{is} x_{jl}$$





$$c_{is,jl} := (d_{ij} - d_{sl})^2$$



Total matching cost: $\sum_{is} c_{is} x_{is} + \sum_{\substack{i \neq j \\ s \neq l}} c_{is,jl} x_{is} x_{jl}$

$$c_{is} := c_{is,is}$$
 - unary costs

$$\min_{x \in \{0,1\}^{n \times n}} \sum_{ijsl} c_{is,jl} x_{is} x_{jl}$$

s.t.:
$$\sum_{s} x_{is} = 1 \ \forall i$$
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$$c_{is,jl} := (d_{ij} - d_{sl})^2$$



Total matching cost: $\sum_{is} c_{is} x_{is} + \sum_{\substack{i \neq j \\ s \neq l}} c_{is,jl} x_{is} x_{jl}$

$$j$$
 $c_{is,jl}$ s

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s.t.:
$$\sum_{s} x_{is} = 1 \ \forall i$$
$$\sum_{i} x_{is} = 1 \ \forall s$$

→ NP-hard

$$c_{is,jl} := (d_{ij} - d_{sl})^2$$



Total matching cost: $\sum_{is} c_{is} x_{is} + \sum_{\substack{i \neq j \\ s \neq l}} c_{is,jl} x_{is} x_{jl}$ $c_{is} := c_{is,is}$ - unary costs $\min_{x \in \{0,1\}^{n \times n}} \sum_{ijsl} c_{is,jl} x_{is} x_{jl}$ $\min_{x \in P} x^\top C x$ s.t.: $\sum x_{is} = 1 \ \forall i$ $\sum x_{is} = 1 \ \forall s$ NP-hard





$$c_{is,jl} := (d_{ij} - d_{sl})^2$$





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"classical" graph matching

Koopmans-Beckmann form



$$c_{is,jl} = \frac{\alpha_{ij}\beta_{sl}}{\beta_{sl}}$$

Lawler form



[Zhou, De La Torre 2012, Factorized Graph Matching]

Factorization into a sum of Koopmans-Beckmann forms

Do we still need optimization in era of NNs?

(except for learning NNs)

Yes, we do:

[Vlastelica et al. '19 Differentiation of blackbox combinatorial solvers]

State-of the art deep graph matching method



Network Architecture

[Rolinek et al. '20 Deep graph matching via black-box differentiation of combinatorial solvers]

Do we still need optimization in era of NNs?

PASCAL VOC Feature points matching:

Method	*	ക്		S i	Ð	₽	Mean	-
GMN-PL	31.1	46.2	58	3.6	83.2	88.6	57.9	
PCA-GM [59]	40.9	55.0	65	7.5	86.7	90.9	63.8	
NGM+ [<mark>60</mark>]	50.8	64.5	59	3.3	81.4	89.6	66.1	
GLMNet [27]	52.0	67.3	63	1.9	79.3	91.3	67.5	LAP Solver
CIE ₁ -H [61]	51.2	69.2	70	5.4	85.2	92.4	68.9	
DGMC* [24]	50.4	67.6	70	9.6	94.3	89.6	73.2 ± 0.5]
BB-GM	61.5	75.0	78	 7.5	97.7	94.4	80.1 ± 0.6	QAP Solver

Everyone uses at least a LAP solver!

[Rolinek et al. '20 Deep graph matching via black-box differentiation of combinatorial solvers]

Bi-Stochastic layer: "Differentiable" linear assignment

$$\min_{x \ge 0} \sum_{i=1}^{n} \sum_{s=1}^{n} c_{is} x_{is} - \rho \mathbf{H}(\mathbf{x})$$

s.t.:
$$\sum_{s=1}^{n} x_{is} = 1 \quad \forall i$$
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Bi-Stochastic layer: "Differentiable" linear assignment

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smooth, differentiable

differentiable Sinkhorn algorithm

Sinkhorn algorithm:

$$\text{Init } x_{is}^0 := \exp(-c_{is}/T)$$

Iterate:

- 1. Normalize each row
- 2. Normalize each column

$$\begin{pmatrix} x_{1,1} & x_{1,2} & \cdots & x_{1,n} \\ x_{2,1} & x_{2,2} & \cdots & x_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ x_{n,1} & x_{n,2} & \cdots & x_{n,n} \end{pmatrix} / \sum_{s=1}^{n} x_{2,s} \\ / \sum_{i=1}^{n} x_{i,2}$$

Operations Research:

• Koopmans-Beckman form (QAPLIB) $c_{is,jl} = \alpha_{ij}\beta_{sl}$

Computer Vision:

Lawler form

Operations Research:

- Koopmans-Beckman form (QAPLIB) $c_{is,jl} = \alpha_{ij}\beta_{sl}$
- C rather dense

- Lawler form
- C often sparse $c_{is} = \infty$ and $c_{is,jl} = 0$

Operations Research:

- Koopmans-Beckman form (QAPLIB) $c_{is,jl} = \alpha_{ij}\beta_{sl}$
- C rather dense
- Precision

- Lawler form
- C often sparse $c_{is} = \infty$ and $c_{is,jl} = 0$
- Speed

Operations Research:

- Koopmans-Beckman form (QAPLIB) $c_{is,jl} = \alpha_{ij}\beta_{sl}$
- C rather dense

- Lawler form
- C often sparse $c_{is} = \infty$ and $c_{is,jl} = 0$

- Precision
- Complete assignment

- Speed
- Incomplete assignment



Incomplete vs. complete assignment



e.g. [Haller et al. 2022. A comparative study of graph matching algorithms in computer vision]















 $[d] = [1, \dots, d] \text{ - set of graphs}$ $C^{[pq]} \text{ - cost matrix for } p \leftrightarrow q \text{ graphs}$ $x^{[pq]}, X^{[pq]} \text{ - } p \leftrightarrow q \text{ assignment}$

Recall QAP: $\min_{x \in P} x^{\top} C x$



 $\begin{bmatrix} d \end{bmatrix} = [1, \dots, d] \text{ - set of graphs} \\ C^{[pq]} \text{ - cost matrix for } p \leftrightarrow q \text{ graphs} \\ x^{[pq]}, X^{[pq]} \text{ - } p \leftrightarrow q \text{ assignment}$

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 $\min_{X \in P} \sum_{p,q \in [d]} (x^{[pq]})^{\top} C^{[pq]} x^{[pq]}$ s.t. $X^{[pq]} X^{[qr]} \leq X^{[pr]} \quad \forall p,q,r \in [d].$

Recall QAP: $\min_{x \in P} x^{\top} C x$



 $[d] = [1, \dots, d] \text{ - set of graphs}$ $C^{[pq]} \text{ - cost matrix for } p \leftrightarrow q \text{ graphs}$ $x^{[pq]}, X^{[pq]} \text{ - } p \leftrightarrow q \text{ assignment}$

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s.t. $X^{[pq]} X^{[qr]} \le X^{[pr]} \quad \forall p,q,r \in [d]$

of [pq] pairs - $O(d^2)$ # of costs $c_{is,jl}^{[pq]}$ - $O(d^2n^4)$ # of cycle consistency constraints - $O(d^3n^3)$

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LAP Solver

Take home messages

Graph matching = matching finite point sets

Linear assignment – point features only

Quadratic assignment – features related to a pair of points

Optimization is a key component



Mean 57.9 63.8

66.1

67.568.9 73.2 ± 0.5

 80.1 ± 0.6

is

