# Tutorial on <br> Inference and Learning in Discrete Graphical Models: Theory and Practice 

ICCV 2015, Santiago de Chile

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December 12th, 2015 (full day)

## About Us



## Schedule

- (08:30-08:40) Opening
- (08:40-09:30) Discrete Graphical Models (50 min)
- Applications in Computer Vision (20 min)
- Definitions and Notation (20 min)
- Overview of Existing Software-packages (10 min)
- (09:30-10:00) Inference in Discrete Graphical Models I (150 min)
- Exact Inference Methods (50 min)
- (10:00-10:30) Coffee Break
- (10:30-12:30) Inference in Discrete Graphical Models II
- ... Exact Inference Methods
- Inference Methods based on Relaxations (40 min)
- Partial Optimality ( 10 min )
- Approximative and Move Making Methods (40 min)
- Meta-Methods : Combining Methods to get a better overall performance ( 10 min )
- (12:15-14:00) Lunch
- (14:00-15:00) From Benchmarks to the Current Limits
- Insights from Benchmark Studies (20 min)
- How to deal with Huge Models and Higher-order Potentials? (20 min)
- Models with Discrete and Continuous Variables (20 min)
- (15:00-15:30) Coffee Break
- (15:30-16:20) Learning in Discrete Graphical Models (50 min)
- Problem Setting
- Maximum Likelihood based Methods
- Prediction-based Parameter Learning Methods
- (16:20-16:30) Closing


## Good to know ...

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- Slides will be online.


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- point You to references with further informations.


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- If You have questions please ask.
- It is more important to get the concepts than the details!


## Applications for Graphical Models in Computer Vision

## Learning and Inference in Graphical Models



## Applications of Graphical Models - outside CV



Label Chemical Structures


VLSI integrated-circuit design


Speech recognition


Flywing partitioning and tracking


Semantic segmentation of the worm

## Applications of Graphical Models - inside CV



Semantic segmentation


Pose Estimation


Depth Estimation



Motion Estimation and Tracking

Debluring and Denoising


## Graphical Models in Computer Vision: A Success Story



- Video Segmentation
- Dense Discrete-Continuous Optimization

Video Enhancement
[Rav-Acha et al. Siggraph 2008]

## Graphical Models in Computer Vision: A Success Story

## Books Sequence Input



- Sparse Graph matching
- 6D Dense Continuous Motion

Scene Flow Estimation
[Abu Alhaija et. al. GCPR 2015]

## Graphical Models in Computer Vision: A Success Story



Input


Output

- Learning Gaussian Markov Random Fields
- State-of-the art deconvolution

Image Deconvolution
[Schmidt et al. CVPR 2013]

## Graphical Models in Computer Vision: A Success Story

Current Leader of Semantic Segmentation Challenge VOC2012

(a) Testing

(b) Truth

(d) Full model

[Efficient Piecewise Training of Deep
Structured Models for Semantic Segmentation, Lin, Shen Reid, Hegel, Arxiv 2015]

## Open Challenges

1. Inference in Large and Complex Models
2. Exact inference
3. Fast Inference
4. Continuous Variables and Mixed Models
5. Deep Learning and Graphical Models

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## 1. Large Scale Models



Drosophila Embryo, Keller Lab, Janelia, US e.g. $2,880 \times 250 \mathrm{MB}$ stack with 100 K cells

Track each individual cell perfectly (99.9\% accuracy needed!)

## 1. Large Scale Models



Fly Wing, Eaton Lab, MPI-CBG, Dresden

Track each individual cell perfectly (99.9 \% accuracy needed!)

## 1. Complex Models

Curvature-based Segmentation [Shekhovtsov et al. GCPR 2010]


Higher-order Potentials penalizing high Curvature


Curvature Model


Pairwise MRF


Curvature Model


Pairwise
MRF

- Right model ... but inference too hard!

1. TRW-S
2. TRW-S (hard constraints)
3. Block-ICM


LB=1.25


## Open Challenges

1. Inference in Large and Complex Models
2. Exact inference
3. Fast Inference
4. Continuous Variables and Mixed Models
5. Deep Learning and Graphical Models

## 2. Exact Inference

## Input: <br> Image sequence


[Data courtesy from Oliver Woodford]

Output: New view


Model: Minimize a binary 4-connected pair-wise graph (choose a colour-mode at each pixel)
[Fitzgibbon et al. ICCV ‘03]

## 2. Exact Inference



Ground Truth


Graph Cut with truncation
[Rother et al. '05]


Belief Propagation (approximate solution)

ICM, Simulated
Annealing



QPBOP
[Boros et al. '06;
Rother et al. '07]
(approximate solution)
(approximate solution) (exact solution)

## Why is the result not perfect? Model or Optimization

## Open Challenges

1. Inference in Large and Complex Models
2. Exact inference
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4. Continuous Variables and Mixed Models
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## 2. Fast Inference



1D cell tracking

## 3. Fast Inference



- Human in-the Loop
- Deep Learning

Joint Segmentation and Tracking [Jug et al. BAMBI (MICCAI) 2014]

## Open Challenges

1. Inference in Large and Complex Models
2. Exact inference
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4. Continuous Variables and Mixed Models
5. Deep Learning and Graphical Models

## 4. Continuous Variables Models



Discrete Variables


Continuous Variables
[Bleyer, Rhemann, Rother. BMVC '2011]

## Open Challenges

1. Inference in Large and Complex Models
2. Exact inference
3. Fast Inference
4. Continuous Variables and Mixed Models
5. Deep Learning and Graphical Models

## 5. Deep Learning and Structured Models

Input


Ground Truth


CNN Trained separately


+ CRF Trained separately (87.6\%)


CNN Trained jointly


+ CRF Trained jointly
(89.0\%)
- CRF gives CNN additional regularization
- What is the optimal combination of CRFs and CNNs?

ML Learning of a generic CNN-CRF model
[Kirrilov et al. arxiv 2015]

## Summary

1. Graphical Models are everywhere in Vision
2. Many interesting open challenges
3. Enjoy the Tutorial!

## Definitions and Notation

## Graphical Models

Definition: Graphical Model
A graphical model is a model for which a graph denotes some structure between variables (represented by its nodes).


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Definition: Graphical Model
A graphical model is a model for which a graph denotes some structure between variables (represented by its nodes).


$$
\begin{aligned}
& \text { (ㄴi. } \Leftrightarrow x_{i} \in X_{i}=\left\{0,1,2, \ldots, L_{i}\right\} \\
& \text { or } \\
& \Omega \subset \mathbb{R}
\end{aligned}
$$

Discrete Variable
Continuous Variable

## Graphical Models: Probabilistic Graphical Model

Definition: Probabilistic Graphical Model
A probabilistic graphical model is a probabilistic model for which a graph $G=(V, E \subset V \times V)$ denotes the conditional independence structure between random variables (represented by its nodes).

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Bayesian Network


$$
\begin{array}{ll}
G \Rightarrow \text { Markov Properties (MP) } & \text {, e.g. } X_{v} \Perp X_{V \backslash \operatorname{de}(v)} \mid X_{\mathrm{pa}(v)} \\
G \Rightarrow \text { Factorization } & , P(X)=\prod_{v \in V} P\left(X_{v} \mid X_{\mathrm{pa}(v)}\right)
\end{array}
$$

$G$ has to be a directed acyclic graph (DAG)

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## Bayesian Network


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$G$ has to be a directed acyclic graph (DAG)

## Markov Random Field (MRF)



$$
G \Rightarrow \text { Markov Properties (MP) } \quad \text {, e.g. }{ }^{1} X_{u} \Perp X_{v} \mid X_{V \backslash\{u, v\}} \Leftrightarrow(u v) \notin E
$$

$\mathrm{MP} \stackrel{*}{\Rightarrow}$ Factorization $\quad, P(X=x) \propto \prod_{C \in \operatorname{cliques}(G)} \varphi_{C}\left(x_{C}\right)$
${ }^{1}$ Pairwise Markov property

* P has to be a positive density or G has to be chordal


## Graphical Models: MRF vs. CRF

## Markov Random Field (MRF)



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## Conditional Random Field (CRF)


$G \Rightarrow$ Markov Properties (MP)

$$
\begin{aligned}
& \text {, e.g. }{ }^{1} X_{u} \Perp X_{v} \mid X_{v \backslash\{u, v\}} \Leftrightarrow(u v) \notin E \\
& , P(X=x \mid D=d) \underset{C \in \operatorname{cliques}(G)}{\propto} \varphi_{C}\left(X_{H(C)}, d_{O(C)}\right)
\end{aligned}
$$

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${ }^{1}$ Pairwise Markov property

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## Graphical Models: Factor Graph Model

## Definition: Factor Graph (Graphical) Model

A factor graph model is a model for which a factor graph
$G=(V, F, E \subset V \times F)$ denotes the factorization of the objective function over its variables (represented by its nodes).

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\text { e.g. } & P(X=x) & =\prod_{f \in F} \varphi_{f}\left(x_{\mathrm{ne}(f)}\right) \\
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$$
\operatorname{func}(x)=\varphi_{f_{1}}\left(x_{n e\left(f_{1}\right)}\right) \otimes \varphi_{f_{2}}\left(x_{n e\left(f_{2}\right)}\right) \otimes \varphi_{f_{3}}\left(x_{n e\left(f_{3}\right)}\right) \otimes \varphi_{f_{4}}\left(x_{n e\left(f_{4}\right)}\right) \otimes \varphi_{f_{5}}\left(x_{n e\left(f_{5}\right)}\right) \quad \otimes \varphi_{f_{6}}\left(x_{n e\left(f_{6}\right)}\right)
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$$
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& =\varphi_{f_{1}}\left(x_{1}\right) \quad \otimes \varphi_{f_{2}}\left(x_{2}, x_{3}\right) \otimes \varphi_{f_{3}}\left(x_{2}, x_{4}\right) \otimes \varphi_{f_{4}}\left(x_{1}, x_{4}\right) \otimes \varphi_{f_{5}}\left(x_{2}, x_{3}, x_{4}\right) \otimes \varphi_{f_{6}}\left(x_{2}, x_{3}, x_{5}, x_{6}\right)
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J(x) \quad & =\varphi_{f_{1}}\left(x_{1}\right) \quad+\varphi_{t_{2}}\left(x_{2}, x_{3}\right)+\varphi_{t_{3}}\left(x_{2}, x_{4}\right)+\varphi_{f_{4}}\left(x_{1}, x_{4}\right)+\varphi_{t_{5}}\left(x_{2}, x_{3}, x_{4}\right) \varphi_{f_{6}}\left(x_{2}, x_{3}, x_{5}, x_{6}\right)
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& J(x) & =\sum_{f \in F} \varphi_{f}\left(x_{\mathrm{ne}(f)}\right)
\end{array}
$$

Definition: Order
The order of a factor is its degree, i.e.

$$
o(f)=|\{v \in V \mid(v, f) \in E\}|
$$

The order of a factor graph is the maximal order of its factors, i.e.

$$
o(G)=\max _{f \in F} o(f)
$$

## Graphical Models

## Conditional Factor Graph Model



## Graphical Models

## Conditional Factor Graph Model



Discrete Factor Graph Model


## Graphical Models

Conditional Factor Graph Model


Discrete Factor Graph Model

## Extended Factor Graph Model


$G=(V, F, E, F, \mathcal{E})$


## Graphical Models

Conditional Factor Graph Model


Extended Factor Graph Model

$G=(V, F, E, \mathcal{F}, \mathcal{E})$

Discrete Factor Graph Model


Weighted Factor Graph Model


## Graphical Models

Conditional Factor Graph Model


Extended Factor Graph Model

$G=(V, F, E, \mathcal{F}, \mathcal{E})$

Discrete Factor Graph Model


Weighted Factor Graph Model

$\boldsymbol{G}=\left(V, F, E, \mathcal{F}, \mathcal{E}, W, E_{W}\right)$

## Graphical Models

Conditional Factor Graph Model


Extended Factor Graph Model

$G=(V, F, E, F, \mathcal{E})$

Discrete Factor Graph Model


Weighted Factor Graph Model

$\boldsymbol{G}=\left(V, F, E, \mathcal{F}, \mathcal{E}, W, E_{W}\right)$

## Graphical Models: Shorthands for Factor Graphs



$$
\begin{aligned}
& G=(V, F, E) \\
& J(x)=\sum_{f \in F} \varphi_{f}\left(X_{\mathrm{ne}(f)}\right)
\end{aligned}
$$

Common used Shorthand for 2nd order Factor Graph Models


$$
\begin{aligned}
& G=(V, E \subset V \times V) \\
& J(x)=\sum_{v \in V} \varphi_{V}\left(x_{V}\right)+\sum_{e \in E} \varphi_{e}\left(x_{e}\right)
\end{aligned}
$$

Important: This is not a MRF, because $G$ decodes directly the factorization!

## Graphical Models: Shorthands for Factor Graphs

Factor Graph Models


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\begin{aligned}
& G=(V, F, E) \\
& J(x)=\sum_{f \in F} \varphi_{f}\left(x_{\mathrm{ne}(f)}\right)
\end{aligned}
$$

Common used Shorthand for general Factor Graph Models


$$
\begin{aligned}
& G=\left(V, E \subset 2^{V}\right) \\
& J(x)=\sum_{v \in V} \varphi_{v}\left(x_{v}\right)+\sum_{C \in E} \varphi_{C}\left(x_{C}\right)
\end{aligned}
$$

Important: This is not a MRF, because $G$ decodes directly the factorization!

## Graphical Models: Factor-Graph Model vs. Markov Random Fields



Factor graph models are more powerful than MRFs/CRFs and Bayesian networks in expressing factorization.

Bayesian networks are more powerful than Factor graph models and MRFs/CRFs in expressing conditional independences.

The terms MRFs/CRFs are often loosely used in computer vision!
Further information: [Lauritzen, 1996, Koller and Friedman, 2009]

## Graphical Models: Gibbs Distribution

- Given an energy function

$$
J(x)=\sum_{C \in \mathcal{C} \subset 2^{v}} \varphi_{C}\left(x_{C}\right)
$$

## Graphical Models: Gibbs Distribution

- Given an energy function

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J(x)=\sum_{C \in \mathcal{C} \subset 2^{v}} \varphi_{C}\left(x_{C}\right)
$$

we can define the Gibbs distribution for a given free parameter $\beta$ by

$$
P(X=x)=p(x)=\frac{1}{Z} \exp (-\beta \cdot J(x))=\frac{1}{Z} \prod_{C \in \mathcal{C} \subset 2^{V}} \exp \left(-\beta \cdot \varphi_{C}\left(x_{C}\right)\right)
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## Graphical Models: Gibbs Distribution

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$$

where the partition function is given by

$$
Z=\sum_{x \in \mathcal{X}} \exp (-\beta \cdot J(x))
$$

## Graphical Models: Gibbs Distribution

- Given an energy function

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$$

where the partition function is given by

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Z=\sum_{x \in \mathcal{X}} \exp (-\beta \cdot J(x)) .
$$

And what is the statistical meaning of $p(x)$ ?


## Graphical Models: Gibbs Distribution

- Given a distribution
$p(x)$ with Markov properties given by $G=(V, E)$


## Graphical Models: Gibbs Distribution

- Given a distribution

$$
p(x) \text { with Markov properties given by } G=(V, E)
$$

- If $p(\cdot)$ is strict positive $(p(x)>0)$ or $G$ is chordal it factorize into

$$
p(x)=\prod_{c \in \mathcal{C} \subset 2^{V}} \phi_{C}\left(x_{C}\right)
$$

where $\mathcal{C}$ the maximal cliques in $G$.
$\rightarrow$ Hammersley-Clifford theorem [Hammersley and Clifford, 1971].

## Graphical Models: Gibbs Distribution

- Given a distribution

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where $\mathcal{C}$ the maximal cliques in $G$.
$\rightarrow$ Hammersley-Clifford theorem [Hammersley and Clifford, 1971].

- We can define a energy function which has $p(x)$ as Gibbs distribution

$$
J(x)=-\frac{1}{\beta} \log (p(x))=\sum_{c \in \mathcal{C} \subset 2^{v}}-\frac{1}{\beta} \log \left(\phi_{C}\left(x_{C}\right)\right)
$$

## Graphical Models: Inference

## Graphical Models: Inference

Probabilistic Inference

$$
\begin{array}{rlr}
p\left(x_{i}\right) & =\sum_{x^{\prime} \in \mathcal{X}, x_{i}^{\prime}=x_{i}} p\left(x^{\prime}\right) & \text { Marginals } \\
p_{\max }\left(x_{i}\right) & =\max _{x^{\prime} \in \mathcal{X}, x_{i}^{\prime}=x_{i}} p\left(x^{\prime}\right) & \text { Max-Marginals } \\
\tilde{x} & \sim p(x) & \text { Sampling } \\
Z & =\sum_{x^{\prime} \in \mathcal{X}} \exp \left(-J\left(x^{\prime}\right)\right) & \text { Partition Function }
\end{array}
$$

MAP Inference and Energy Minimization

```
\(x^{*} \in \arg \max p(x)\)
    \(x \in \mathcal{X}\)
\(x^{*} \in \arg \min J(x) \quad\) Configuration with the Lowest Energy
```


## Graphical Models: Inference

Using MAP-inference to approximate marginals

- Perturb and MAP
- Expected MAP solution of perturbed model $\Rightarrow$ marginals
- [Papandreou and Yuille, 2011, Hazan et al., 2013]
- Frank-Wolf Algorithm
- Perturbed MAP-problems show up in Frank-Wolf algorithm for calculating marginals
- [R. Krishnan, 2015]


## Graphical Models: Inference

MAP Inference and Energy Minimization

$$
\begin{array}{lr}
x^{*} \in \underset{x \in \mathcal{X}}{\arg \max } p(x \mid d) & \text { Most Likely Explanation } \\
x^{*} \in \underset{x \in \mathcal{X}}{\arg \min } J(x \mid d) & \text { Configuration with the Lowest Energy }
\end{array}
$$



## Graphical Models: Inference

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\begin{array}{lr}
x^{*} \in \underset{x \in \mathcal{X}}{\arg \max } p(x \mid d) & \text { Most Likely Explanation } \\
x^{*} \in \underset{x \in \mathcal{X}}{\arg \min } J(x \mid d) & \text { Configuration with the Lowest Energy }
\end{array}
$$



## Graphical Models: Inference

MAP Inference and Energy Minimization

$$
\begin{array}{lr}
x^{*} \in \underset{x \in \mathcal{X}}{\arg \max } p(x \mid d) & \text { Most Likely Explanation } \\
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## Graphical Models: Learning

## Graphical Models: Learning

## Graphical Models: Learning

1. Select (Learn) the (number of) variables and labels and their meaning

(3)

## Graphical Models: Learning

1. Select (Learn) the (number of) variables and labels and their meaning
2. Select (Learn) the structure of the model


## Graphical Models: Learning

1. Select (Learn) the (number of) variables and labels and their meaning
2. Select (Learn) the structure of the model
3. Learn independently local potentials/feature-functions

- physical priors
- handcrafted features
- output of random forest or any local model from pattern recognition textbook
- output of CNNs


$$
J(x)=\sum_{i=1}^{3} \varphi_{f_{i}}\left(x_{i}\right)+\varphi_{f_{12}}\left(x_{1}, x_{2}\right)+\varphi_{f_{23}^{1}}\left(x_{2}, x_{3}\right)+\varphi_{f_{23}^{2}}\left(x_{2}, x_{3}\right)
$$

## Graphical Models: Learning

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4. Learn global parameters that adjust local terms (optional)


$$
J(x)=\sum_{i=1}^{3} \varphi_{f_{i}}\left(x_{i}\right)+\varphi_{f_{12}}\left(x_{1}, x_{2}\right)+\varphi_{f_{23}}\left(x_{2}, x_{3}\right)+\varphi_{f_{23}^{2}}\left(x_{2}, x_{3}\right)
$$

$$
J(x)=\sum_{i=1}^{3} w_{f_{i}}{\overline{f_{i}}}\left(x_{i}\right)+w_{f_{12}} \bar{\varphi}_{f_{12}}\left(x_{1}, x_{2}\right)+w_{f_{12}^{\prime}} \bar{\varphi}_{f_{23}^{\prime}}\left(x_{2}, x_{3}\right)+w_{f_{12}^{2}} \bar{\varphi}_{f_{23}^{2}}\left(x_{2}, x_{3}\right)
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## Graphical Models: Learning

1. Select (Learn) the (number of) variables and labels and their meaning
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\begin{gathered}
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\end{gathered}
$$

## Graphical Models: Learning

## Problem Setting

Given some training-data $\left(d^{1}, \ldots, d^{N}\right)$ and ground truth configuration ( $x^{G T ; 1}, \ldots, x^{G T ; N}$ ) we would like to learn the optimal model parameters.

## Graphical Models: Learning

## Problem Setting

Given some training-data $\left(d^{1}, \ldots, d^{N}\right)$ and ground truth configuration ( $x^{G T ; 1}, \ldots, x^{G T ; N}$ ) we would like to learn the optimal model parameters.

Maximum Likelihood Learning

$$
\left.w^{*} \in \arg \max _{w} \prod_{i} p\left(x^{G T ; i} \mid w, d^{i}\right)\right)
$$

Maximum Margin Learning

$$
w^{*} \in \arg \min _{w} \sum_{i} \Delta\left(x^{G T ; i}, \arg \max _{x \in X} p\left(x \mid w, d^{i}\right)\right)+\lambda\|w\|_{2}^{2}
$$

## Software Packages

## Software Packages

| Library | Authors | Language | Last Updated | License | Model | Inference | Learning |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
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Models: discrete factor graph models (DFG), Bayesian Nets (BN), ${ }^{P}$ only pairwise, ${ }^{S}$ restiction on the graph structure Models: probabilistic inference ( $P$ ), energy minimization ( $\mathbf{E}$ )
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## Which Library is the Best? Depends on Your Goal!

functionality generality / flexibility

## How to use OpenGM



## How to use OpenGM



## How to use OpenGM



## Inference Methods for Energy Minimization with Discrete Graphical Models

## Energy Minimization Methods available in OpenGM



## Exact Inference Methods for Energy Minimization with Discrete Graphical Models

## Exact Inference Methods



## Exact Inference Methods for Energy Minimization with Discrete Graphical Models

$$
\arg \min _{x \in \mathcal{X}} \sum_{f \in F} \varphi_{f}\left(x_{\text {ne(f) }}\right)
$$

## Exact Inference Methods for Energy Minimization with Discrete

 Graphical Models$$
\arg \min _{x \in \mathcal{X}} \sum_{f \in F} \varphi_{f}\left(x_{\text {ne(f) }}\right)
$$

Stop! This problem is NP-hard, so exact inference is in general not tractable!

## Exact Inference Methods for Energy Minimization with Discrete Graphical Models

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\arg \min _{x \in \mathcal{X}} \sum_{f \in F} \varphi_{f}\left(x_{\text {ne(f) }}\right)
$$

Stop! This problem is NP-hard, so exact inference is in general not tractable!


|  | Problem Restriction | Runtime Complexity |
| :--- | :--- | :--- |
| Dynamic Programming |  |  |
| $\rightarrow$ Viterbi Algorithm | acyclic structure | polynomial |
| $\rightarrow$ Junction Tree Algorithm | limited tree-width | polynomial in the tree-width <br> polynomial in the model order per iteration <br> Reduction to |
| $\rightarrow$ Max Flow | no, but not exact |  |
| $\rightarrow$ Submodular Minimization | pairwise submodular | polynomial |
| $\rightarrow$ Perfect Matching | submodular | polynomial |
| $\rightarrow$ Multicut / Multiway Cut Problem | (outer) planar binary | polynomial |
| Redts-(like) functions | exponential in the worst case |  |
| $\rightarrow$ Brute Force Search Problem |  | no |
| $\rightarrow$ Best First Search | no | exponential |
| Integer Linear Program | no | exponential in the worst case |

[^0]
## Exact Inference by Dynamic Programming (aka Viterbi algorithm)

## Exact Inference Methods: Dynamic Programming on a Chain

Optimization Problem


$$
\arg \min _{x_{1}, x_{2}, x_{3}} \varphi_{f_{1}}\left(x_{1}\right)+\varphi_{t_{2}}\left(x_{2}\right)+\varphi_{t_{3}}\left(x_{3}\right)+\varphi_{t_{12}}\left(x_{1}, x_{2}\right)+\varphi_{f_{23}}\left(x_{2}, x_{3}\right)
$$

## Exact Inference Methods: Dynamic Programming on a Chain

Optimization Problem


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\arg \min _{x_{1}, x_{2}, x_{3}} \varphi_{f_{1}}\left(x_{1}\right)+\varphi_{f_{2}}\left(x_{2}\right)+\varphi_{f_{3}}\left(x_{3}\right)+\varphi_{f_{12}}\left(x_{1}, x_{2}\right)+\varphi_{f_{23}}\left(x_{2}, x_{3}\right)
$$

Optimization by Dynamic Programming


$$
\min _{x_{1}} \varphi_{f_{1}}\left(x_{1}\right)+\min _{x_{2}} \varphi_{f_{12}}\left(x_{1}, x_{2}\right)+\varphi_{t_{2}}\left(x_{2}\right)+\min _{x_{3}} \varphi_{f_{23}}\left(x_{2}, x_{3}\right)+\varphi_{f_{3}}\left(x_{3}\right)
$$

## Exact Inference Methods: Dynamic Programming on a Chain

Optimization Problem


$$
\arg \min _{x_{1}, x_{2}, x_{3}} \varphi_{f_{1}}\left(x_{1}\right)+\varphi_{f_{2}}\left(x_{2}\right)+\varphi_{f_{3}}\left(x_{3}\right)+\varphi_{f_{12}}\left(x_{1}, x_{2}\right)+\varphi_{f_{23}}\left(x_{2}, x_{3}\right)
$$

Optimization by Dynamic Programming


$$
\min _{x_{1}} \varphi_{f_{1}}\left(x_{1}\right)+\min _{x_{2}} \varphi_{f_{12}}\left(x_{1}, x_{2}\right)+\varphi_{f_{2}}\left(x_{2}\right)+\varphi_{\bar{F}_{2}}\left(x_{2}\right)
$$

## Exact Inference Methods: Dynamic Programming on a Chain

Optimization Problem


$$
\arg \min _{x_{1}, x_{2}, x_{3}} \varphi_{f_{1}}\left(x_{1}\right)+\varphi_{t_{2}}\left(x_{2}\right)+\varphi_{t_{3}}\left(x_{3}\right)+\varphi_{t_{12}}\left(x_{1}, x_{2}\right)+\varphi_{f_{23}}\left(x_{2}, x_{3}\right)
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Optimization by Dynamic Programming


## Exact Inference Methods: Dynamic Programming on a Chain

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## Exact Inference Methods: Dynamic Programming on a Chain

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Optimization by Dynamic Programming


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$$

Optimization by Dynamic Programming


## Exact Inference Methods: Dynamic Programming on a Tree

Factor Graph Model


## Exact Inference Methods: Dynamic Programming on a Tree

Factor Graph Model


## Exact Inference Methods: Dynamic Programming on a Tree

 Factor Graph Model

Optimization by Dynamic Programming


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Optimization by Dynamic Programming


## Exact Inference Methods: Dynamic Programming on a Tree

 Factor Graph Model

Optimization by Dynamic Programming


## Exact Inference Methods: Dynamic Programming on a Higher-Order Tree

Factor Graph Model


Algorithmic Idea


## Exact Inference Methods: Dynamic Programming on General Graphs



## Exact Inference Methods: Dynamic Programming on General Graphs



Why did't You merge node 2 and 3

## Exact Inference Methods: Dynamic Programming on General Graphs



Why did't You merge node 2 and 3

## Exact Inference Methods: Dynamic Programming on General Graphs



Why did't You merge node 2 and 3 ... and node 4 and $5 \ldots$

## Exact Inference Methods: Dynamic Programming on General Graphs



Why did't You merge node 2 and 3 ... and node 4 and $5 \ldots$

## Exact Inference Methods: Dynamic Programming on General Graphs



Why did't You merge node 2 and 3
... and node 4 and $5 \ldots$ now it looks

like a tree!

## Exact Inference Methods: Dynamic Programming on General Graphs

## Background

- Idea: Reduce inference to dynamic programming on the Junction Tree
- The junction tree is a cluster tree that fulfills the Running Intersection Property. It exists if and only if the node-adjacency graph is chordal.
- Finding the triangulation that generates a graph with the lowest clique-size is NP-hard.

[Lauritzen, 1996, Koller and Friedman, 2009]


## (Loopy) Belief Propagation (Not always exact, but most often useful)

## Belief Propagation



## Belief Propagation



## Messages

$$
\begin{aligned}
& m_{v \rightarrow f}\left(x_{v}\right)=\sum_{f^{\prime} \in \operatorname{ne}(v) \backslash\{f\}} m_{f^{\prime} \rightarrow v}\left(x_{v}\right) \\
& m_{f \rightarrow v}\left(x_{v}\right)=\min _{x_{\mathrm{ne}(f) \backslash\{v\}} \in X_{\mathrm{ne}(f) \backslash\{v\}}} \varphi_{f}\left(x_{\mathrm{ne}(f)}\right)+\sum_{u \in \operatorname{ne}(f) \backslash\{v\}} m_{u \rightarrow f}\left(x_{v}\right)
\end{aligned}
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## Belief Propagation



## Messages

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## Belief Propagation



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& m_{f \rightarrow v}\left(x_{v}\right)=\sum_{x_{\text {ne }(f) \backslash\{v\}} \in X_{\text {ne }(f) \backslash\{v\}}} \varphi_{f}\left(x_{\text {ne }(f)}\right)+\sum_{u \in \operatorname{ne}(f) \backslash\{v\}} m_{u \rightarrow f}\left(x_{v}\right)
\end{aligned}
$$

## Damping

$$
m^{\text {new }}\left(x_{v}\right)=(1-\alpha) \cdot m\left(x_{v}\right)+\alpha \cdot m^{\text {old }}\left(x_{v}\right)
$$

## Belief Propagation



## Messages

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Normalization

## Belief Propagation



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$$

Works surprisingly good for cyclic models...

## Belief Propagation

Message Schedule

- Parallel $\rightarrow$ Loopy Belief Propagation (LBP) [Pearl, 1988]
- Sequential $\rightarrow$ Sequential Belief Propagation (BPS) [Felzenszwalb and Huttenlocher, 2006]
- Informed $\rightarrow$ Residual Belief Propagation (RBP)
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Message Update Rules (from a theoretical point of view)
- Original Updates: Optimize a non-convex objective function $\rightarrow$ lack of convergence
- Modified Updates: Optimize a convex objective function $\rightarrow$ we will come back to this point later


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[3.6 [Kschischang et al., 2001] $\rightarrow$ generalization to other semi-rings


## Reduction to Min-st-Cut/Max-Flow Problem (aka GraphCut)

## Minimal st-Cut Problem

## Definition: Min-st-Cut Problem

Given a weighted directed graph $G=(V, E, w)$ with a source-node $s \in V$ and a sink-node $t \in V$. Find the subset of edges $E^{\prime} \subset E$ with the minimal edge-weight $\sum_{e \in E^{\prime}} w(e)$ such that no path from $s$ to $t$ exists in $G^{\prime}=\left(V, E \backslash E^{\prime}\right)$.


## Remark

If the edge-weights are non-negative the Min-st-Cut problem is efficiently solvable.

## Max Flow Problem

## Definition: Max-Flow Problem

Given a weighted directed graph $G=(V, E, w)$ with a source-node $s \in V$ and a sink-node $t \in V$. Find the maximal flow passing from $s$ to $t$, where a positive flow smaller than $w(e)$ flow can be passed over the edge $e$ and conservation holds.


## Max-Flow Min-Cut Theorem

The maximum value of an st flow with capacity $w(e)$ is equal to the minimum st cuts with edge weights $w(e)$.

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## Exact Inference Methods: Reformulate into Min-st-Cut

Requirements:

- Binary label-space
- Regular/submodular pairwise terms

Reformulation:


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Requirements:

- Binary label-space
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Reformulation:


## Reduction to Submodular Minimization

## Submodular Minimization

$$
\arg \min _{S \in 2^{V}} f(S)
$$

Submodular Function
If $V$ is a finite set, a submodular function is a set function $f: 2^{V} \rightarrow \mathbb{R}$, where $2^{V}$ denotes the power set of $V$, for which for every $S, T \subseteq V$ the inequality $f(S)+f(T) \geq f(S \cup T)+f(S \cap T)$ holds.

## Minimizing Submodular Function

Finding the subset $S \subset V$ that minimize a submodular function is computable in polynomial time [Schrijver, 2000, lyer et al., 2013]

Relation to Binary Graphical Models
Let $S$ be the set of variables taking label 0 and $V \backslash S$ the set of labels taking label 1.

ICML-2013 Tutorial by Stefanie Jegelka and Andreas Krause
(http://submodularity.org/)

## Reduction to Perfect Matching via Max-Cut on Planar Graphs

## Max Cut Problem

## Definition:

Given a weighted undirected graph $G=(V, E, w)$ the Max-Cut problem is to find a cut $E^{\prime} \subset E$, that divide $V$ into two sets, that the sum of the weights of cut edges is maximized.


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## Definition:

Given a weighted undirected graph $G=(V, E, w)$ the Max-Cut problem is to find a cut $E^{\prime} \subset E$, that divide $V$ into two sets, that the sum of the weights of cut edges is maximized.


Remark:
For planar graphs the Max Cut problem can be efficiently
solved [Kasteleyn, 1961, Fisher, 1961] [Globerson and Jaakkola, 2006,
Schraudolph and Kamenetsky, 2008] by a reduction to a Perfect Matching problem, e.g. Blossom V [Kolmogorov, 2009], on some expanded dual graph.

## Reduction to a Max Cut Problem

$$
x_{i} \in\{0,1\} \quad \varphi_{f_{i j}}\left(x_{i}, x_{j}\right)=-\lambda_{i j} \mathbb{I}\left(x_{i} \neq x_{j}\right) \quad \varphi_{u_{i}}\left(x_{i}\right)=-\alpha_{i} \mathbb{I}\left(x_{i}=0\right)
$$

## Planar Potts Model without Unaries


$x_{i}=$ cluster number of node $i$

Outer-Planar Potts Model with Unaries

$x_{i}= \begin{cases}1 & \text { if the edge }(t i) \text { is not cut } \\ 0 & \text { otherwise }\end{cases}$
(30) [Schraudolph and Kamenetsky, 2008]

## Reduction to Minimal Multicut Problem

## Minimal Multicut Problem

Definition:
Given a weighted undirected graph $G=(V, E, w)$ the Minimal Multicut problem is to find a cut $E^{\prime} \subset E$, that divide $V$ into a unknown number of sets, that the sum of the weights of cut edges is minimized.


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Given a weighted undirected graph $G=(V, E, w)$ the Minimal Multicut problem is to find a cut $E^{\prime} \subset E$, that divide $V$ into a unknown number of sets, that the sum of the weights of cut edges is minimized.


Remark:
The above multicut is not a cut, since the clustering is not 2-colorable!

## Reduction to a Minimal Multicut Problem

$$
x_{i} \in\{0, \ldots, L\} \quad \varphi_{t_{i j}}\left(x_{i}, x_{j}\right)=\lambda_{i j} \mathbb{I}\left(x_{i} \neq x_{j}\right) \quad \varphi_{u_{i}}\left(x_{i}\right)=\alpha_{i ; x_{i}}
$$

Potts Model with $L=|V|$ and without Unaries $\rightarrow$ Multicut

$x_{i}=$ cluster number of node $i$

Potts Model with Unaries $\rightarrow$ Multiway Cut


$$
\begin{aligned}
& \left(\begin{array}{cccc}
0 & 1 & \cdots & 1 \\
1 & \ddots & \ddots & \vdots \\
\vdots & \ddots & \ddots & 1 \\
1 & \cdots & 1 & 0
\end{array}\right)\left(\begin{array}{c}
w_{t_{i} i} \\
\vdots \\
\vdots \\
w_{t_{L} i}
\end{array}\right)=\left(\begin{array}{c}
\alpha_{i ; 1}-\alpha_{i ; 0} \\
\vdots \\
\vdots \\
\alpha_{i ; L}-\alpha_{i ; 0}
\end{array}\right) \\
& x_{i}= \begin{cases}l & \text { if the edge }\left(t_{i} i\right) \text { is not cut } \\
0 & \text { otherwise }\end{cases}
\end{aligned}
$$

## Reduction to Multicut/Multiway Cut Problem

```
Reduction to Max Flow
```

Label-space: Unary terms: Pairwise terms:
Problem structure:
Runtime complexity:
$x_{i} \in\{0,1\}$ arbitrary
$\varphi_{f_{i j}}\left(x_{i}, x_{j}\right)=\lambda_{i j} \mathbb{I}\left(x_{i} \neq x_{j}\right) \quad \lambda_{i j} \in \mathbb{R}^{+} \quad$ (submodular) arbitrary
polynomial

## Reduction to Multicut/Multiway Cut Problem

| Reduction to Max Flow |  |  |  |
| :---: | :---: | :---: | :---: |
| Label-space: | $x_{i} \in\{0,1\}$ |  |  |
| Unary terms: | arbitrary |  |  |
| Pairwise terms: | $\varphi_{f_{i j}}\left(x_{i}, x_{j}\right)=\lambda_{i j} \mathbb{I}\left(x_{i} \neq x_{j}\right)$ | $\lambda_{i j} \in \mathbb{R}^{+}$ | (submodular) |
| Problem structure: | arbitrary |  |  |
| Runtime complexity: | polynomial |  |  |
| Reduction to Max Cut |  |  |  |
| Label-space: | $x_{i} \in\{0,1\}$ |  |  |
| Unary terms: | none ${ }^{1}$ or arbitrary ${ }^{2}$ |  |  |
| Pairwise terms: | $\varphi_{t_{i j}}\left(x_{i}, x_{j}\right)=\lambda_{i j} \mathbb{I}\left(x_{i} \neq x_{j}\right)$ | $\lambda_{i j} \in \mathbb{R}$ |  |
| Problem structure: | planar ${ }^{1}$ or outer planar ${ }^{2}$ |  |  |
| Runtime complexity: | polynomial |  |  |

## Reduction to Multicut/Multiway Cut Problem



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planar ${ }^{1}$ or outer planar ${ }^{2}$
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Reduction to Multicut / Multiway Cut
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Pairwise terms:
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Runtime complexity:
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```
```

arbitrary
exponential in the worst case, but often tractable in practice

## Solving a Multicut/Multiway Cut Problem by an (I)LP

LP Formulation

$$
\begin{aligned}
& \min _{y \in[0,1]^{E}} \sum_{e \in E} w_{e} \cdot y_{e} \\
& \text { s.t. } A y \leq b
\end{aligned}
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This is no valid multicut! $\Rightarrow A y \leq b$

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s.t. $A y \leq b$
exponential size
Idea of Cutting-Plane Methods


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Idea of Cutting-Plane Methods


- Separation Procedure:

Find violated constraints

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Idea of Cutting-Plane Methods


- Separation Procedure: Find violated constraints
- The separation procedure is as hard as the original problem or the original problem is as hard as the separation procedure
(8) Ellipsoid method for combinatorial optimization [Grötschel et al., 1981]


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## ILP Formulation

$$
\begin{aligned}
\min _{y \in[0,1]^{E}} & \sum_{e \in E} w_{e} \cdot y_{e} \\
\text { s.t. } & \tilde{A} y \leq \tilde{b} \\
& y \in\{0,1\}^{E}
\end{aligned}
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## Reduction to Shortest Path Search

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$\min _{x \in X} \sum_{f \in F} \varphi_{f}(x) \Leftrightarrow$ Finding shortest path from $s$ to $t$

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## Reduction to Shortest Path Search

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Useful Ideas

## Reduction to Shortest Path Search

edge weights


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## MAP Inference as Integer Linear Program

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$$
\begin{array}{ll} 
& \min _{\mu}\langle\theta, \mu\rangle \\
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## MAP inference as Integer Linear Program

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\begin{array}{ll}
\min _{\mu}\langle\theta, \mu\rangle \\
\text { s.t. } & A \mu \leq b \\
& \mu \in\{0,1\}^{N}
\end{array}
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\end{aligned}
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## MAP inference as Integer Linear Program

$$
\min _{\mu}\langle\theta, \mu\rangle
$$

s.t. $\quad A \mu \leq b$

$$
\mu \in\{0,1\}^{N}
$$

Fine, but in general ILPs cannot be solved in polynomial time!

$$
\forall v \in V:
$$

$$
\sum_{x_{v} \in X_{v}} \mu_{v ; x_{c}}=1
$$

$$
\stackrel{\mu \in\{0,1\}^{N}}{\Rightarrow} \sum_{x_{v} \in X_{v}} \mathbb{I}\left(\mu_{v ; x_{v}}=1\right)=1
$$

$$
\forall f \in F, v \in n e(f):
$$

$$
\sum_{x_{n e(f) \backslash v} \in X_{n e(f) \backslash v}} \mu_{f ; x_{n e(f)}}=\mu_{v ; x_{v}}
$$

$$
\mu \stackrel{\{0,1\}^{N}}{\Rightarrow} \mu_{f ; x_{n e(f)}}=\prod_{v \in n e(f)} \mu_{v, x_{v}}
$$

## MAP inference as Integer Linear Program

Reasons why ILPs had been ignored

- Worst case complexity is exponential
- ILPs are often very memory consuming
- Good ILP-solvers are expensive


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## Reasons why ILPs are Relevant

- Worst case complexity does not always matter!
- Highly optimized commercial solvers (e.g. CPLEX, Gurobi) are freely available for academic use.
- ILPs can compute the global optimal solution.
- No limitations on the model . . if we ignore memory requirements.
- Always a good baseline for small problems.


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- Always a good baseline for small problems.

How ILP-solvers works

- Branch and Bound
- Searching for generic constraints
- The black magic is to combine all of this


## Finally some Simple Tricks ...

## Tricks to make Inference more Tractable

Ring


## Tricks to make Inference more Tractable

Ring


- Solve the acyclic problem for each possible labeling of $x_{8}$.
- Select the best solution from all these problems.


## Tricks to make Inference more Tractable

Partially Acyclic Graph


## Tricks to make Inference more Tractable

Partially Acyclic Graph


- Presolve acyclic substructures by dynamic programming.
- Solve core problem, with acyclic part replaced by unary factor
- Calculate labeling for acyclic part as for dynamic programming given the solution of the core problem.
(Kappes et al., 2013]


## Tricks to make Inference more Tractable

Permuted Submodular


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Permuted Submodular

not regular $\xrightarrow{?}$ not submodular

## Tricks to make Inference more Tractable

Permuted Submodular

not regular $\xrightarrow{?}$ not submodular

regular $\rightarrow$ submodular

- The reformulation is variant to label permutation $\rightarrow$ the order of labels matters!
- Reorder labels
- Solve max-flow or submodular-minimization problem
- Undo reorder for solution
[Schlesinger, 2007, Swoboda et al., 2013]


## Inference Methods based on Relaxations



## MAP Inference as Integer Linear Program

$$
(\arg ) \min _{x \in \mathcal{X}} \sum_{v \in V} \theta_{v}\left(x_{v}\right)+\sum_{u v \in E} \theta_{u v}\left(x_{u}, x_{v}\right)
$$



## MAP Inference as Integer Linear Program



## MAP Inference as Integer Linear Program

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\begin{array}{r}
(\arg ) \min _{x \in \mathcal{X}} \sum_{v \in V} \theta_{v}\left(x_{v}\right)+\sum_{u v \in E} \theta_{u v}\left(x_{u}, x_{v}\right)=(\arg ) \min _{x \in \mathcal{X}}\langle\theta, \delta(x)\rangle \\
\delta(x)=\underbrace{100}_{u} \underbrace{010000000}_{u v} \underbrace{010}_{v}
\end{array}
$$


(arg) $\min _{x \in \mathcal{X}}\langle\theta, \delta(x)\rangle$

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$(\arg ) \min _{\mu \in \operatorname{conv}(\mathcal{X})}\langle\theta, \mu\rangle$

## MAP Inference as Integer Linear Program

$$
(\arg ) \min _{x \in \mathcal{X}} \sum_{v \in V} \theta_{v}\left(x_{v}\right)+\sum_{u v \in E} \theta_{u v}\left(x_{u}, x_{v}\right)=(\arg ) \min _{x \in \mathcal{X}}\langle\theta, \delta(x)\rangle
$$



$$
\begin{gathered}
(\arg ) \min _{\mu \in \operatorname{conv}(\delta(\mathcal{X}))}\langle\theta, \mu\rangle \\
\operatorname{conv}(\delta(\mathcal{X})) \Rightarrow A \mu \geq b
\end{gathered}
$$

## Relaxed MAP Inference



$$
\begin{aligned}
& \min _{\mu}\langle\theta, \mu\rangle \\
\text { s.t. } \mu= & \sum_{x \in \mathcal{X}} \alpha_{x} \delta(x) \\
& \sum_{x \in \mathcal{X}} \alpha_{x}=1 \\
& \alpha_{x} \geq 0, x \in \mathcal{X}
\end{aligned}
$$

## Relaxed MAP Inference



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Exponentially many constraints!

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s.t. $A \mu \geq b, A$ is small!

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& \alpha_{x} \geq 0, x \in \mathcal{X}
\end{aligned}
$$

Exponentially many constraints!
$\mu$ - non-relaxed solution

$\Rightarrow$
$\min _{\mu}\langle\theta, \mu\rangle$
s.t. $A \mu \geq b, A$ is small!

$$
\begin{aligned}
& \forall v \in V: \quad \sum_{x_{v} \in X_{v}} \mu_{v ; x_{c}}=1 \\
& \forall f \in F, v \in n e(f): \\
& \quad \sum_{x_{\text {ne(f) }} \in X_{\text {ne(f)\v}}} \mu_{f ; x_{n e(f)}}=\mu_{v ; x_{v}} \\
& \quad \mu \in\{0,1\}^{N} \mu \in[0,1]^{N}
\end{aligned}
$$

$\Leftarrow \quad \mu$ - relaxed solution

## Relaxed MAP Inference



## Local Polytope Complexity



| $\mid$ ㄴ\| | $\|X\|$ | vertices in $\mathbb{L}(M)$ |
| ---: | ---: | ---: |
| 2 | 8 | 12 |
| 3 | 27 | 207 |
| 4 | 64 | 8.992 |
| 5 | 125 | 853.725 |

## Why dedicated solvers needed for the relaxed inference?



Pascal VOC 2012 semantic segmentation model $\approx 500 \times 300 \times 21$ labels

$2 \cdot 10^{6}$ variables
$>10^{9}$ variables

## Why dedicated solvers needed for the relaxed inference?



Pascal VOC 2012 semantic segmentation model $\approx 500 \times 300 \times 21$ labels

$2 \cdot 10^{6}$ variables
$>10^{9}$ variables

Standard LP solvers (simplex, interior point) do not scale well!

## Lagrangean (Dual) Decomposition

Certain problems are easily solvable (e.g. acyclic with dynamic programming):

Optimization by Dynamic Programming


$$
\min _{x_{1}} \varphi_{f_{1}}\left(x_{1}\right)+\min _{x_{2}} \varphi_{f_{12}}\left(x_{1}, x_{2}\right)+\varphi_{t_{2}}\left(x_{2}\right)+\varphi_{F_{2}}\left(x_{2}\right)
$$

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$$

What about decomposing the problem into solvable subproblems?

## Lagrangean (Dual) Decomposition



## Lagrangean (Dual) Decomposition



$$
\theta_{v}\left(x_{v}\right) \quad=\quad \frac{\theta_{v}\left(x_{v}\right)}{2}+\lambda_{v}\left(x_{v}\right)
$$

$+$
$E\left(\theta^{r}, x\right)$


## Lagrangean (Dual) Decomposition



$$
\theta_{v}\left(x_{v}\right)
$$

$\min _{x \in \mathcal{X}} E(\theta, x)$
$+$
$+$

$$
=
$$

$$
\begin{aligned}
& = \\
& =
\end{aligned}
$$


$=\quad \frac{\theta_{v}\left(x_{v}\right)}{2}+\lambda_{v}\left(x_{v}\right)$
$+$
$\frac{\theta_{v}\left(x_{v}\right)}{2}-\lambda_{v}\left(x_{v}\right)$


$$
E\left(\theta^{r}, x\right)
$$


$\stackrel{H \lambda}{2}$
$+$

## Lagrangean (Dual) Decomposition



## Lower Bound Optimization

$$
\begin{gathered}
\theta^{c}=\frac{\theta_{v}}{2}+\lambda_{v} ; \quad \theta^{r}=\frac{\theta_{v}}{2}-\lambda_{v} \\
\max _{\lambda}\left[\min _{x^{c} \in \mathcal{X}} E\left(\theta^{c}(\lambda), x^{c}\right)+\min _{x^{r} \in \mathcal{X}} E\left(\theta^{r}(\lambda), x^{r}\right)\right]
\end{gathered}
$$

## Lower Bound Optimization

$$
\theta^{c}=\frac{\theta_{v}}{2}+\lambda_{v} ; \quad \theta^{r}=\frac{\theta_{v}}{2}-\lambda_{v}
$$



## Lower Bound Optimization



## Lower Bound Optimization



## Lower Bound Optimization



Optimization:

- convex
- non-smooth
- large-scale


## First-Order Convex Optimization

| update rule | subproblem | rate | note |
| :---: | :---: | :---: | :---: |
| sub-gradient | MAP-inference | $O\left(\frac{1}{\epsilon^{2}}\right)$ | step-size selection |
| mirror-descent | MAP-inference | $O\left(\frac{1}{\epsilon^{2}}\right)$ |  |
| bundle | MAP-inference | $O\left(\frac{1}{\epsilon^{2}}\right)$ | additional QP |
| coord. ascent | min-marginals | unknown | no optimum guarantee |
| smooth coord. ascent | probab.-marginals | unknown | exp. operation |
| smooth acc. ascent | marginalization | $O\left(\frac{1}{\epsilon}\right)$ | exp. operation |
| proximal (e.g. ADMM) | proximal inference | $O\left(\frac{1}{\epsilon}\right)$ |  |

## First-Order Convex Optimization: Subgradient Method



## First-Order Convex Optimization: Subgradient Method



## First-Order Convex Optimization: Subgradient Method



First-Order Convex Optimization: Subgradient Method

$$
\begin{aligned}
D(\lambda)=\min _{x^{c}}\left\langle\frac{\theta}{2}+\lambda, \delta\left(x^{c}\right)\right\rangle+\min _{x^{r}}\left\langle\frac{\theta}{2}-\lambda,\right. & \left.\delta\left(x^{r}\right)\right\rangle \\
& \Rightarrow \partial D(\lambda)=\delta\left(x^{* c}\right)-\delta\left(x^{* r}\right)
\end{aligned}
$$



First-Order Convex Optimization: Subgradient Method

gradient $\nabla D$
Gradient ascent:

$$
\begin{gathered}
\lambda^{t+1}=\lambda^{t}+\tau \nabla D \\
L \geq \tau>0
\end{gathered}
$$

$$
D(\lambda)=\min _{x^{c}}\left\langle\frac{\theta}{2}+\lambda, \delta\left(x^{c}\right)\right\rangle+\min _{x^{r}}\left\langle\frac{\theta}{2}-\lambda, \delta\left(x^{r}\right)\right\rangle
$$

$$
\Rightarrow \partial D(\lambda)=\delta\left(x^{* c}\right)-\delta\left(x^{* r}\right)
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+ based on MAP inference

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$$

+ based on MAP inference
- sensible to $\tau^{t}$


## Subgradient method:

$\lambda^{t+1}=\lambda^{t}+\tau^{t} \partial D$
$\tau^{t}>0, \tau^{t} \rightarrow 0, \sum_{t=1}^{\infty} \tau^{t}=\infty$

$$
\Rightarrow \partial D(\lambda)=\delta\left(x^{* c}\right)-\delta\left(x^{* r}\right)
$$

First-Order Convex Optimization: Subgradient Method

gradient $\nabla D$
Gradient ascent:

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$$

$$
\Rightarrow \partial D(\lambda)=\delta\left(x^{* c}\right)-\delta\left(x^{* r}\right)
$$

+ based on MAP inference
- sensible to $\tau^{t}$
- slow: converges as $O\left(\frac{1}{\epsilon^{2}}\right)$

$$
\epsilon=0.1 \Rightarrow t=100
$$

$$
\epsilon=0.01 \Rightarrow t=10000
$$

## First-Order Convex Optimization: Bundle Method



## First-Order Convex Optimization: Bundle Method

Bundle:

$\lambda^{t+1}=\arg \min _{\lambda} \widehat{D}(\lambda)+\gamma^{t}\left\|\lambda-\lambda^{t}\right\|^{2}$

$$
\gamma^{t} \rightarrow 0
$$

Subgradient:


$$
\begin{gathered}
\lambda^{t+1}=\lambda^{t}+\tau^{t} \partial D \\
\tau^{t} \rightarrow 0
\end{gathered}
$$

## First-Order Convex Optimization: Bundle Method

Bundle:
$\lambda^{t+1}=\arg \min _{\lambda} \widehat{D}(\lambda)+\gamma^{t}\left\|\lambda-\lambda^{t}\right\|^{2}$
$\gamma^{t} \rightarrow 0$
(30appes et al., 2012]


$$
\begin{gathered}
\lambda^{t+1}=\lambda^{t}+\tau^{t} \partial D \\
\tau^{t} \rightarrow 0
\end{gathered}
$$

(30) [Storvik and Dahl, 2000,

Schlesinger and Giginyak, 2007,
Komodakis et al., 2011]

+ based on MAP inference
+ less sensible to $\gamma^{t}$
+ much faster convergence in practice
- can be slow: worst-case complexity $O\left(\frac{1}{\epsilon^{2}}\right)$


## First-Order Convex Optimization: Coordinate Ascent

$$
\begin{array}{r}
\lambda_{i}^{t+1}=\arg \min _{\lambda_{i}} D\left(\lambda_{1}^{t}, \ldots, \lambda_{i}, \ldots, \lambda_{N}^{t}\right) \\
i=1, \ldots, N
\end{array}
$$

## First-Order Convex Optimization: Coordinate Ascent

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## First-Order Convex Optimization: Coordinate Ascent

$\lambda_{i}^{t+1}=\arg \min _{\lambda_{i}} D\left(\lambda_{1}^{t}, \ldots, \lambda_{i}, \ldots, \lambda_{N}^{t}\right)$

$$
i=1, \ldots, N
$$

+ requires MAP - inference (min-marginals)
+ can be very efficiently implemented
- can get stuck
- convergence rate is unknown


Inference algorithms:
TRW-S: [Kolmogorov, 2006]
SRMP: [Kolmogorov, 2015]
Max-Sum-Diffusion (MSD): [Schlesinger and Antoniuk, 2011, Werner, 2007]
MPLP: [Globerson and Jaakkola, 2007]
Norm-Product BP (NPBP): [Hazan and Shashua, 2010]

## First-Order Convex Optimization: Smoothing



## First-Order Convex Optimization: Smoothing



$$
\begin{aligned}
& \min \quad \rightarrow \text { "soft" min } \\
& \min _{x^{c}} E^{c}\left(x^{c}\right) \rightarrow-T \log \sum_{x^{c}} \exp \left(-E^{c}\left(x^{c}\right) / T\right) \\
& \text { MAP-inf. } \rightarrow \text { Probabilistic inf. }
\end{aligned}
$$

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+ accelerated gradient ascent converges as $O\left(\frac{1}{\epsilon}\right)$

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- Expensive 'exp' and 'log' operations


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$$

+ coordinate ascent does not gut stuck!
+/- Requires probabilistic inference: tractable for acyclic subproblems
- Expensive 'exp' and 'log' operations

Used to approximate probabilistic inference on the master graph.

Smoothing theory: [Nesterov, 2005] Inference algorithms:
Nesterov: [Savchynskyy et al., 2011]
ADSal: [Savchynskyy et al., 2012]
See also: [Johnson et al., 2007, Werner, 2009, Meshi et al., 2012]

## Rounding, Obtaining a Labeling



## Rounding, Obtaining a Labeling



Column- or row-wise labeling as a solution

## Rounding, Obtaining a Labeling



Column- or row-wise labeling as a solution

## Rounding, Obtaining a Labeling



Column- or row-wise labeling as a solution


## Rounding: Alternatives



How to round?

- Best of integer solutions, collected over iterations

- Local rounding: deterministic/probabilistic: $\mu \geq 0.5$ ?
[Ravikumar et al., 2010, Kleinberg and Tardos, 2002]
- Sequential conditional rounding (TRW-S) [Kolmogorov, 2006]


## Rounding: Alternatives



How to round?

- Best of integer solutions, collected over iterations

- Local rounding: deterministic/probabilistic: $\mu \geq 0.5$ ?
[Ravikumar et al., 2010, Kleinberg and Tardos, 2002]
- Sequential conditional rounding (TRW-S)
[Kolmogorov, 2006]

There is no single best method. Most of methods - heuristics.

## Tree Agreement



## Tree Agreement



Strong tree agreement in each node $=$ solution of the non-relaxed problem


## Tree Agreement



In general: only weak tree agreement holds in the limit


- How to obtain the primal relaxed solution:[Savchynskyy and Schmidt, 2014]


## Acyclic Subgraphs $\Rightarrow$ Local Polytope Relaxation

Acyclic subgraphs for decomposition:

|  | subgraph | problematic solvers |
| :---: | :---: | :---: |
|  | 1-edge graph | no |
|  | chain | proximal inference |

## Acyclic Subgraphs $\Rightarrow$ Local Polytope Relaxation

Acyclic subgraphs for decomposition:

|  | subgraph | problematic solvers |
| :---: | :---: | :---: |
|  | 1-edge graph | no |
|  | chain | proximal inference |

Theorem ([Komodakis et al., 2011])
Arbitrary covering acyclic subgraphs $\Rightarrow$ the same local polytope relaxation.

Which decomposition is better? Smooth. coord. ascent, grid models


Synthetic uniformly generated $200 \times 100$ grid, 5 labels

Which decomposition is better? Smooth. coord. ascent, grid models


Color segmentation, Potts model, $360 \times 240$ grid, 12 labels

Which decomposition is better? Smooth. coord. ascent, grid models


Stereo reconstruction (tsukuba), $384 \times 288$ grid, 16 labels

Which decomposition is better? Smooth. coord. ascent, grid models


Bigger subproblems $\Rightarrow$ less iterations!


## Which update rule is better? Stereo, Columns+Rows decomposition



## Which update rule is better? Stereo, Columns+Rows decomposition




## Local Polytope: Algorithm's Overview



## Local Polytope: Algorithm's Overview



## Beyond Acyclic Subproblems


(3trandmark and Kahl, 2010]

## Beyond Acyclic Subproblems


(306) [Kappes et al., 2010]

## Beyond Acyclic Subproblems




## Beyond Acyclic Subproblems


(a) 1-fan

(b) 4-fan

(c) 5-fan


Hint: Subgraphs, where LP relaxation is loose, are preferable as subproblems.

## Partial Optimality

## Partial Optimality: Definition





## Partial Optimality: Definition



LP relaxation

- Is the integer part of the relaxed solution optimal?
- Can we eliminate labels with zero weight?

In general - no, but sometimes - yes.

## 2-Label Local Polytope, QPBO

## Recall:



Min-st-Cut
(30) [Boros and Hammer, 2002, Rother et al., 2007a]

## 2-Label Local Polytope, QPBO




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## What about $\geq 3$ labels?

a) MQPBO

MQPBO-method (submodular relaxation) [Kohli et al., 2008]:

- Convert $n$-label to 2-label problem. ("Battleship" encoding, [Schlesinger and Flach, 2006] )
- Apply QPBO.



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- Apply QPBO.


Polytope of the submodular relaxation $\supseteq$ local polytope Depends on the selected order of variables [Swoboda et al., 2013] In practice works up to 3-4 variables only

## What about $\geq 3$ labels?

b) Arbitrary relaxation




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LP relaxation


- Is the integer part of the relaxed solution optimal?



## Criterion:

Check whether for all $y$ the labeling $x$ remains optimal: $\forall y: x=\arg \min _{x^{\prime}} E\left(x^{\prime}, y\right)$
integer $x$ and fractional $y$ labelings

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- relaxed inference $\arg \min _{x^{\prime}} E\left(x^{\prime}, y\right)$ is sufficient;
- efficient procedure exists, for local polytope and a more general question.


## Partial Optimality:Potts Models



Figure : [Shekhovtsov et al., 2015], Color segmentation, Potts model (> 20 instances)

## Progress in Partial Optimality



## Alternative to Partial Optimality: CombiLP

Partial optimality:

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CombiLP: [Savchynskyy et al., 2013]


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Check whether for optimal $y^{*}$ the labeling $x$ remains optimal: $x=\arg \min _{x^{\prime}} E\left(x^{\prime}, y^{*}\right)$
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## Alternative to Partial Optimality: CombiLP <br> Partial optimality:



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+ Weaker requirement $\Rightarrow$ stronger result.


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integer $x$ and fractional $y$ labelings

+ Weaker requirement $\Rightarrow$ stronger result.
- CombiLP has exponential complexity (Partial optimality - polynomial).


## CombiLP vs. Partial Optimality

Partial optimality: [Shekhovtsov et al., 2015]


CombiLP: [Savchynskyy et al., 2013]


## Approximative and Move Making Methods

## Move Making Methods



## Move Making Methods:

- Start from any solution $X$
- Maintain current best feasible solution $\hat{X}$
- Try to improve $\hat{X}$ with a set of Moves
- Stop when no more improvement is possible with the set of Moves
- Work in the primal domain

Advantages:

- trivial warm start
- can improve solutions from arbitrary solvers
- fast and scalable


## Downsides:

- (often) no lower bound / guarantees
- can get stuck in local minim
- "hope" for a good local minimum


## Subgraph Methods

## Iterated Conditional Modes [Besag, 1986]



- Start from arbitrary starting point (e.g. $X=0$ )


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(11) and (16) are not connected!
They can be optimized independent!

SITIUILameousty

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- Optimize ${ }^{\sqrt{11}}$, $\sqrt{1 / 2}$ and $\sqrt{\sqrt{5}}$ simultaneously
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## Subgraph Construction



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## Lazy Flipper [Andres et al., 2010]

- systematically optimize all connected subgraphs of size $k$
- for $k=|V|$ global optimal
- for $k=1$ equal to ICM [Besag, 1986]


Connected Subgraph Tree:
Maximum Subgraph size $k=6$


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Connected Subgraph Tree:
Maximum Subgraph size $k=5$


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Connected Subgraph Tree:
Maximum Subgraph size $k=4$


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Connected Subgraph Tree:
Maximum Subgraph size $k=3$


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Connected Subgraph Tree:
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## Local Rules for Global MAP: When Do They Work ? [Jung et al., 2009]

Local, iterative randomized PTAS for MAP:


- Optimize "ball"-shaped subgraphs around random nodes
- radius of "ball" drawn from truncated geometric distribution $Q$

Image Credit:[Jung et al., 2009]

## Local Rules for Global MAP: When Do They Work? [Jung et al., 2009]

Local, iterative randomized PTAS for MAP:
Graph G


- Optimize "ball"-shaped subgraphs around random nodes
- radius of "ball" drawn from truncated geometric distribution $Q$
Image Credit: [Jung et al., 2009]
Provable $\epsilon$ - approximation:
- for certain geometric distribution $Q$
- on graphs with polynomial growth rate
- with $n \log ^{2} n$ iterations
- pairwise MRFS


## Fusion Move Methods

## Fusion Moves[Lempitsky et al., 2010]

- Energy function $J(X)$ with large label space $|X|$

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J(X)=\sum \phi_{f}\left(x_{n e(f)}\right)
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- Binary Label Space:

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\hat{X}=\left\{x \in X \mid \forall i: x_{i} \in\left\{x^{1}, x^{2}\right\}\right\}
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- Allowed Moves:
$X^{\text {Move }}=\left\{x \in X \mid J(x) \leq \min \left(J\left(x^{1}\right), J\left(x^{2}\right)\right)\right\}$


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Multi-Label Energy Function:


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## Fusion Move Based Algorithms



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[Kolmogorov and Zabih, 2002]

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Swap:[Kolmogorov and Zabih, 2002]

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- Inference Based Proposals



## Fusion Moves for Correlation Clustering [Beier et al., 2015]

$$
J(X)=\sum_{u v \in E} \quad \omega_{u v} \cdot x_{u} \neq x_{v} \quad\left|X_{i}\right|=|V|
$$

(a) current best $y^{\prime}$

(c) $\hat{y}=y^{\prime} \cup y^{\prime \prime}$
(d) contracted graph $G_{\hat{y}}$
(e) CC on $G_{\hat{y}}$
(b) proposal $y^{\prime \prime}$

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(e) CC on $G_{\hat{y}}$
(f) result on $G$

## Local Submodular Approximations

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Non Submodular Energy $J(X)$ :

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\begin{gathered}
J(X)=\sum \varphi_{i}\left(x_{i}\right)+\sum \varphi_{i j}\left(x_{i}, x_{j}\right) \quad x_{i} \in\{0,1\} \\
\tilde{J}^{x^{0}}(x) \approx J(X) ; \quad \tilde{J}^{x^{0}}(x) \geq J(x) ; \quad \tilde{J}^{x^{0}}\left(x_{0}\right)=J\left(x_{0}\right)
\end{gathered}
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- Compute submodular approximation $\tilde{J}^{x^{0}}(x)$ arround current labeling
- LSA-TR and LSA-AUX [Gorelick et al., 2014a]


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- Compute submodular approximation $\tilde{J}^{x^{2}}(x)$ arround current labeling
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$$
\begin{gathered}
J(X)=\sum \varphi_{i}\left(x_{i}\right)+\sum \varphi_{i j}\left(x_{i}, x_{j}\right) \quad x_{i} \in\{0,1\} \\
\tilde{\jmath}^{x^{0}}(x) \approx J(X) ; \quad \tilde{J}^{x^{0}}(x) \geq J(x) ; \quad \tilde{J}^{x^{0}}\left(x_{0}\right)=J\left(x_{0}\right)
\end{gathered}
$$



- Compute submodular approximation $\tilde{J}^{x^{2}}(x)$ arround current labeling
- Approximation $\tilde{J^{2}}(x)$ is only valid close to current labeling
- By optimizing $\tilde{J}$ the next solution is generated
- LSA-TR and LSA-AUX [Gorelick et al., 2014a]


## Polyhedral Interpretation



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- Each move can be interpreted as an optimization of an inner polytope


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## Meta-Methods: Combining Methods to get better overall Performance

## Fusion Moves With Inference Based Proposals



- Some algorithms do not decrease the energy monotonously


## Fusion Moves With Inference Based Proposals



- Some algorithms do not decrease the energy monotonously
- Make monotone by remembering current best solution


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- Some algorithms do not decrease the energy monotonously
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- Used generated labels more efficient


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- Investigated in detail by Kappes and Beier[Kappes et al., 2014] with TRWS Dual Decomposition.


## Fusion Moves With Inference Based Proposals



- Some algorithms do not decrease the energy monotonously
- Make monotone by remembering current best solution
- Used generated labels more efficient
- Compute Fusion move between $>0$ and $y^{1}$
- Proposed by Lempitsky[Lempitsky et al., 2010] with Loopy BP
- Investigated in detail by Kappes and Beier[Kappes et al., 2014] with TRWS Dual Decomposition.
- OpenGM allows this trick for all inference algorithms


## Fusion Moves For Parallelization

## Alg1

Alg2

## Alg3

## Alg 4

- Run many different algorithms / proposal generators in parallel


## Fusion Moves For Parallelization



- Run many different algorithms / proposal generators in parallel
- Hierarchically fuse them


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Multi Scale Methods


Multi Scale Methods[Bagon and Galun, 2012]


## Multi Scale Methods

- build energy pyramid

- [Bagon and Galun, 2012]


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## Multi Scale Methods

- build energy pyramid
- Optimize top down

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- build energy pyramid
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- Warm start with solution from layer above
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## Multi Scale Methods



## Multi Scale Methods

- build energy pyramid
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- Warm start with solution from layer above
- [Bagon and Galun, 2012]
- [Meir et al., 2015]


## Insights from Benchmark Studies

## Benchmarks for Graphical Models

- Middlebury MRF [Szeliskiet al., 2008]
- Probabilistic Inference Challenge 2011 http://www.cs.huji.ac.il/project/PASCAL/
- OpenGM Benchmark [Kappes et al., 2015]


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```
http://www.cs.huji.ac.il/project/PASCAL/
```

- OpenGM Benchmark [Kappes et al., 2015]

Evaluations from other communities:

- MAX-CSP 2008 Competition
- Max-SAT Evaluation(s)

```
http://maxsat.ia.udl.cat/introduction/
```


## OpenGM Datasets: Overview


(a) Pixel-based Models

(c) Unsupervised Partitioning

(b) Superpixel-based Models

(d) Higher-order Models

(f) Small but hard models

(e) Large-scale Models

## ... 32 datasets, over 2000 problem instances




## Typical Result Table

| algorithm | runtime | value | bound | mem |  |  | accuracy |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\alpha$-Exp-QPBO | 0.01 sec | -866.85 | - $-\infty$ | 0.01 | 587 | 0 | 0.7694 |
| ogm-LBP-LF2 | 0.06 sec | -866.76 | - $-\infty$ | 0.01 | 576 | 0 | 0.7699 |
| ogm-LF-3 | 0.45 sec | -866.27 | $-\infty$ | 0.01 | 420 | 0 | 0.7699 |
| ogm-TRWS-LF2 | 0.01 sec | -866.93 | $-866.93$ | 0.01 | 714 | 712 | 0.7693 |
| BPS-TAB | 0.10 sec | -866.73 | $-\infty$ | 0.01 | 566 | 0 | 0.7701 |
| ogm-BPS | 0.02 sec | -866.77 | $-\infty$ | 0.01 | 585 | 0 | 0.7694 |
| ogm-LBP-0.95 | 0.02 sec | -866.76 | - $-\infty$ | 0.01 | 580 | 0 | 0.7696 |
| ogm-TRBP-0.95 | 0.11 sec | -866.84 | - $-\infty$ | 0.01 | 644 | 0 | 0.7708 |
| ogm-TRBPS | 0.13 sec | -866.79 | - $-\infty$ | 0.01 | 644 | 0 | 0.7705 |
| ADDD | 0.06 sec | -866.92 | $-866.93$ | 0.01 | 701 | 697 | 0.7693 |
| MPLP | 0.04 sec | -866.91 | $-866.93$ | 0.01 | 700 | 561 | 0.7693 |
| MPLP-C | 0.04 sec | -866.92 | -866.93 | 0.01 | 710 | 567 | 0.7693 |
| ogm-ADSAL | 0.04 sec | -866.93 | -866.93 | 0.01 | 714 | 712 | 0.7693 |
| ogm-BUNDLE-H | 0.26 sec | $-866.93$ | $-866.93$ | 0.01 | 715 | 673 | 0.7693 |
| ogm-BUNDLE-A+ | 0.07 sec | $-866.93$ | -866.93 | 0.01 | 715 | 712 | 0.7693 |
| ogm-LP-LP | 0.23 sec | -866.92 | -866.93 | 0.05 | 712 | 712 | 0.7693 |
| TRWS-TAB | 0.01 sec | $-866.93$ | $-866.93$ | 0.01 | 714 | 712 | 0.7693 |
| BRAOBB-1 | 17.61 sec | $-866.90$ | $-\infty$ | 0.27 | 670 | 0 | 0.7688 |
| ADDD-BB | 0.11 sec | $-866.93$ | $-866.93$ | 0.01 | 715 | 715 | 0.7693 |
| ogm-CombiLP | 0.02 sec | $-866.93$ | $-866.93$ | 0.03 | 715 | 715 | 0.7693 |
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How to select the best method for my problem (without checking all methods)?

## Inference Requirements



## Model Characterization

## Model Characterization:

Size:

- \# variables
- \# factors
- \# labels
- model order



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Size:

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Reduction possibility (partial optimality):

- pairwise Potts
- binary pairwise - QPBO
- binary higher order
- pairwise with tight LP relaxation
- low tree-width sub-graphs
 (junction-tree algorithm)


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Reduction possibility (partial optimality):

- pairwise Potts
- binary pairwise - QPBO
- binary higher order
- pairwise with tight LP relaxation
- low tree-width sub-graphs
 (junction-tree algorithm)
- Tightness of LP relaxation.


## Inference Method Selection



## LP-Relaxation Based Methods: tight?



## Combinatorial Methods:



Move Making Methods: | $\begin{array}{c}\text { Move making } \\ \text { works fine? }\end{array}$ |
| :---: |



O FastPD is typically much faster than $\alpha$-exp

## Inference Method Selection

OpenGM2
Library＊Contacts • Algorithms－References • Benchmarks：cvPR 2013 Arxiv 2014 iJcV 2015 more＊

## OpenGM Benchmark（IJCV 2015）

Benchmark database of discrete energy minimization problems．For further details see
Jörg H．Kappes，Bjoem Andres，Fred A．Hamprecht，Christoph Schnörr，Sebastian Nowozin，Dhruv Batra，Sungwoong Kim，Thorben Kroeger，Bernhard X． Kausler，Jan Lellmann，Bogdan Savchynskyy，Nikos Komodakis，Carsten Rother：
＂A Comparative Study of Modern Inference Techniques for Discrete Energy Minimization Problems＂
International Journal of Computer Vision 2015.
［publisher］，［preprint］，［sublementary material 1］［sublementary material 2］，［bib］

## Download Code and Data

－Optimization methods provided within OpenGM 2．3．3
－Links to download the models are given below．
Models

|  | In－Painting（N4） <br> J．Lellmann et．al． <br> comerted by I．Leifmann and／H．Kappes | 目国图 | Variables | Labels | Order | Stru | Func | Ins | Reference Comment |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | 14400 | 4 | 2 | grid4 | potts | 2 | ［57］ |
|  |  |  |  |  |  |  |  |  |  |
|  | In－Painting（N8） <br> J．Lellmann et．al． <br> converted by／Lellmann and JH．Kappes | 目国 | 14400 | 4 | 2 | grid8 | potts | 2 | ［57］ |
|  | ColorSegmentation（N4） <br> J．Lellmann et．al． <br> converted by I．Lellmann and／H．Kappes | 目国圆 | 76800 | 3.12 | 2 | grid4 | potts | 2 | ［57］ |
| $1$ | ColorSegmentation（N8） <br> J．Lellmann et．al． <br> corverted by／．Lellmann and／H．Kapoes | 目閶 | 76800 | 3.12 | 2 | grids | potts | 2 | ［57］ |

## We are looking for your data and code！

## Huge Models and Higher-Order Potentials?

## What make problems hard?

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1. Inherent Combinatorial Complexity

- LP-relaxations are not tight
- Local optimal decisions do not lead to global optimal decisions.


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2. Higher-Order Factors

- A factor of order $N$ has $L^{N}$ entries, which all have to be explored when no additional information is given.


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- Memory requirements can quickly become very huge


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4. Huge Label-Spaces

- A second-order factor with 10.000 states per variable has $10^{8}$ entries, which all has to be explored when no additional information is given.


## Inherent Combinatorial Complexity

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## When does (real) combinatorial problem show up?

## Inherent Combinatorial Complexity

When does (real) combinatorial problem show up?

- When factors/functions encode constraints
- When we learn a model with many parameters $\rightarrow$ over-fitting
$\rightarrow$ When a set of factors/functions conflicting (frustrated cycles)


## Inherent Combinatorial Complexity

When does (real) combinatorial problem show up?

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$\rightarrow$ When a set of factors/functions conflicting (frustrated cycles)

Examples

- Clustering (consistency constraint) [Andres et al., 2011]
- Graph Matching with weak local assignments (1-to-1 constraint) [Torresani et al., 2008, Komodakis and Paragios, 2008]
- Decision Tree Fields [Nowozin et al., 2011]
- Vector Compression [Babenko and Lempitsky, 2014]


## Inherent Combinatorial Complexity

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## Ways Out

1. Problem specific Constraints/Separation [Nowozin and Lampert, 2009, Kappes et al., 2015b]
2. Make moves over "meaningful" subsets [Gorelick et al., 2014b, Kappes et al., 2014]
3. Stronger local terms can make the problems easier
4. Try to separate the hard combinatorial parts of the problem [Kappes et al., 2013, Kappes et al., 2015b]

## Higher-Order Factors

## Higher-Order Factors

Complexity of Inference in a Factor Graph Model

$$
\Theta\left(\max _{f \in F} \prod_{u \in \operatorname{ne}(f)}\left|X_{u}\right|\right) \approx \Theta\left(L^{o(G)}\right)
$$

where $L$ is the number of labels and $o(G)$ the order of the model

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$$
O\left(L^{2}\right)
$$

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Yes, by using structure of functions.


## Reformulation of Higher order Factors

## General Way to Reformation in LPs

- $\prod_{v \in \text { ne( } f \text { ) }}\left|X_{v}\right|$ slack variables ( $\sim L^{0}$ )
- $\sum_{v \in \text { ne( }(f)}\left|X_{v}\right|$ constraints of size $\left|X_{\text {ne(f) }(f) \backslash v}\right|\left(\sim L \cdot o \times L^{0-1}\right)$
$L=$ number of labels, $o=$ order of factor


## Reformulation of Higher order Factors

General Way to Reformation in LPs

- $\prod_{v \in \text { ne }(f)}\left|X_{v}\right|$ slack variables ( $\sim L^{\circ}$ )
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Idea
Replace a higher-order term by some slack variables + additional constraints.

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Idea
Replace a higher-order term by some slack variables + additional constraints.
Example: Sparse Function [Rother et al., 2009, Kappes et al., 2015a]

$$
\varphi\left(x_{1,2,3,4}\right)= \begin{cases}\gamma & \text { if } x_{1,2,3,4}=(2,6,2,4) \\ 0 & \text { else }\end{cases}
$$

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\begin{aligned}
& \varphi\left(x_{1,2,3,4}\right)=\left\{\begin{array}{lll}
\gamma & \text { if } x_{1,2,3,4} & =(2,6,2,4) \\
0 & \text { else } & \\
& s \leq \mathbb{I}\left(x_{1}=2\right) \\
& & s \leq \mathbb{I}\left(x_{2}=6\right)
\end{array}\right. \\
& \min _{s} \gamma \cdot s \quad \text { s.t. } \quad s \leq \mathbb{I}\left(x_{3}=2\right) \\
& \\
& s \leq \mathbb{I}\left(x_{4}=4\right) \\
& s \geq \mathbb{I}\left(x_{1}=2\right)+\mathbb{I}\left(x_{2}=6\right)+\mathbb{I}\left(x_{3}=2\right)+\mathbb{I}\left(x_{4}=4\right)-3
\end{aligned}
$$

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& \\
=\min _{s} \gamma \cdot s & \text { s.t. } \\
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& s \leq \mathbb{I}\left(x_{3}=2\right) \\
& s \geq \mathbb{I}\left(x_{4}=4\right) \\
& \\
& \\
&
\end{aligned}
$$

## Reformulation of Higher order Factors

General Way to Reformation in LPs

- $\prod_{v \in \text { ne(f) }}\left|X_{v}\right|$ slack variables ( $\sim L^{\circ}$ )
- $\sum_{v \in \operatorname{ne}(f)}\left|X_{v}\right|$ constraints of size $\left|X_{\mathrm{ne}(f) \backslash v \mid}\right|\left(\sim L \cdot o \times L^{0-1}\right)$
$L=$ number of labels, $o=$ order of factor

Idea
Replace a higher-order term by some slack variables + additional constraints.
Examples: Label Cost [Delong et al., 2012]

$$
\begin{aligned}
\varphi^{\prime}\left(x_{V}\right) & = \begin{cases}\gamma_{1} & \text { if } \exists v \in V: x_{V}=1 \\
0 & \text { otherwise }\end{cases} \\
\varphi\left(x_{V}\right) & =\sum_{l \in L} \varphi^{\prime}\left(x_{V}\right)
\end{aligned}
$$

## Reformulation of Higher order Factors

General Way to Reformation in LPs

- $\prod_{v \in \text { ne }(f)}\left|X_{v}\right|$ slack variables ( $\sim L^{\circ}$ )
- $\sum_{v \in \text { ne(f) }}\left|X_{v}\right|$ constraints of size $\left|X_{\text {ne(f) }(f)}\right|\left(\sim L \cdot o \times L^{0-1}\right)$
$L=$ number of labels, $o=$ order of factor

Idea
Replace a higher-order term by some slack variables + additional constraints.
Examples: Label Cost [Delong et al., 2012]

$$
\begin{aligned}
\varphi^{\prime}\left(x_{V}\right) & = \begin{cases}\gamma_{1} & \text { if } \exists v \in V: x_{V}=1 \\
0 & \text { otherwise }\end{cases} \\
\varphi\left(x_{V}\right) & =\sum_{l \in L} \varphi^{\prime}\left(x_{V}\right)
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$$


requires that $\gamma_{l} \geq 0$

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requires that $\gamma_{l} \geq 0$

For $|L|=2$ the label costs $c_{l}$ can be also negative

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$P^{N}$ Potts [Kohli et al., 2009]

$$
\varphi_{f}\left(x_{n e(f)}\right)= \begin{cases}\gamma_{k} & \text { if } x_{v}=k, \forall v \in \operatorname{ne}(f) \\ 0 & \text { otherwise }\end{cases}
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requires that $\gamma_{k} \leq 0, \forall k$

## Order Reduction Higher-Order Boolean Functions

$$
\varphi:\{0,1\}^{N} \rightarrow \mathbb{R}
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Reducing Negative-Coefficient Monomials [Freedman and Drineas, 2005]

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\varphi(x)=-x_{1} \cdots x_{N}=\min _{y \in\{0,1\}} y\left((N-1)-\sum_{i=1}^{N} x_{i}\right)
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\begin{aligned}
\varphi(x) & =x_{1} \cdots x_{N} \\
& = \begin{cases}\min _{y \in\{0,1\}^{N}} \sum_{i=1}^{\lfloor N-1 / 2\rfloor} y_{i}\left(1\left(-\sum_{j=1}^{N} x_{j}+2 i\right)\right)+\sum_{i<j} x_{i} x_{j} & \text { if } \bmod (N, 2)=1 \\
\min _{y \in\{0,1\}^{N}} \sum_{i=1}^{\lfloor N-1 / 2\rfloor} y_{i}\left(2\left(-\sum_{j=1}^{N} x_{j}+2 i\right)\right)+\sum_{i<j} x_{i} x_{j} & \text { if } \bmod (N, 2)=0\end{cases}
\end{aligned}
$$

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\end{aligned}
$$

(Ishikawa, 2009, Fix et al., 2011, Gallagher et al., 2011]

## Huge Number of Variables

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## Problems

- Model does not fit in the memory
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## Possible Ways Out

- Superpixels/Supervariables [Kim et al., 2011]
- Domain-Decomposition (Dual Decomposition) [Schwing et al., 2011, Strandmark and Kahl, 2010]

- Block-ICM


## Huge Label-Spaces

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How to deal with huge label-spaces?

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- Efficient updates [Felzenszwalb and Huttenlocher, 2006]

$$
m\left(x_{1}\right)=\min _{x_{2}}\left|x_{1}-x_{2}\right|+g\left(x_{2}\right)
$$

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$$
x_{i} \in\{1, \ldots, 10\}^{3}
$$

$$
f_{12}\left(x_{1}, x_{2}\right)=\sum_{i=1}^{3} f_{12}^{i}\left(x_{1}, x_{2^{i}}\right)
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## Mixed Models with Continuous Variables

## Possible Solvers

- Block ICM (most popular)
- Particle-based Methods
- Others: Discrete-Continuous Fusion Move; Gradient Descent with relaxed discrete Variables; Variable elimination ...


## Examples in Computer Vision

## Many Discrete, Few Continuous Variables



Input
Segmentation and Human Pose fitting [Kohli, Rihan, Bray, Torr, IJCV 2008]


Segmentation and Color Model fitting GrabCut [Rother, Kolmogorov, Blake 2004]

## Illustration Block-ICM



## Examples in Computer Vision

## Many Discrete, Many Continuous Variables



Joint Models [Vineet, Rother, Torr, NIPS 2013]

## Examples in Computer Vision

## Many Continuous Variables


time $=1$

time $=2$


Motion

## Stereo Matching



Local stereo matching: check photo-consistency

## Stereo Matching



Local stereo matching: check photo-consistency

## None Front-to-Parallel Surface



## None Front-to-Parallel Surface



3 continuous parameters (depth + normal) for each pixel

## Motivation

We now show a comparison of our slanted window algorithm with the competitors described in the paper.

## PatchMatch Stereo

1. Random initialization
2. Go through pixels in sequential order:
3. Consider solution from left/top neighbour
4. Sample around current solution

1D example:



Left image


Left and right disparity maps (intermediate step of iteration 1)

## PatchMatch Stereo

## The Reindeer Pair

## Why does it work?

- Random Initialization is in your favour


Left image


Ground truth disparities
 $\sim 80.000$ guesses for yellow plane

## What's Missing



PatchMatch Stereo result

## Continuous Variable MRF

Add pairwise terms:


## Pairwise term

$$
\begin{gathered}
\psi_{s t}\left(\boldsymbol{u}_{s}, \boldsymbol{u}_{t}\right)=\omega\left[\boldsymbol{u}_{s} \neq \boldsymbol{u}_{t}\right] \\
\beta w_{s t}\left(\left|\mathbf{n}_{s} \cdot\left(\mathbf{x}_{t}-\mathbf{x}_{s}\right)\right|+\left|\mathbf{n}_{t} \cdot\left(\mathbf{x}_{s}-\mathbf{x}_{t}\right)\right|\right)
\end{gathered}
$$



Cost $=0$
both planes are aligned in 3D


Cost $\neq 0$ :
local curvature or discontinuity

## Solvers

- w/o Pairwise Terms: PatchMatch
- w/ Pairwise Terms (super high-dimensional u):
- Gradient descent + Fusion move
- Simulated Annealing
- Continuous Belief Propagation, e.g. Particle BP


## Discrete State BP [Pearl ‘88]



Sequential schedule

## Algorithm

1. Initialize Messages
2. Go over all Messages
3. Update Message $M_{t \rightarrow s}\left(\boldsymbol{u}_{s}\right)$
4. Compute final Output $\boldsymbol{u}_{s}{ }^{*}=\operatorname{argmin} \mathrm{B}_{s}\left(\boldsymbol{u}_{s}\right)$

Step 3: Update Message


Step 4: Compute neg-log Belief


## Toy Example - Shift 4.0



Target

## Toy Example - Shift 4.0



## Toy Example - Shift 4.0



Error: 0.618; Unary only


Ground Truth

$12 \times 12$ discrete labels

## Toy Example - Shift 4.2



## Toy Example - Shift 4.2



Ground Truth

## Toy Example - Shift 4.2


$12 \times 12$ discrete labels

Error: 0.66

## Toy Example - Shift 4.5



Ground Truth

$12 \times 12$ discrete labels

## Max-Product Particle BP

Each pixel has different set of particles:


Step 3: Move Particles $\left\{\boldsymbol{u}_{s}\right\}$


Noise

## Toy Example - Shift 4.5




Ground Truth


## PatchMatch BP

[Besse, Rother, Fitzgibbon, Kautz, BMVC '12]
Each pixel has different set of particles:


Sequential schedule

$\mathrm{B}_{\mathrm{t}}\left(\boldsymbol{u}_{t}\right)$, i.e. neg. log Belief
s)


Algorithm

1. Initialize Particles and Messages
2. Go over all Messages
3. Move and add Particles $\left\{\boldsymbol{u}_{s}\right\}$
4. Update Message $M_{t \rightarrow s}\left(\boldsymbol{u}_{s}\right)$
5. Take $K$ best Particles wrt Belief $B_{S}\left(\boldsymbol{u}_{s}\right)$
6. Compute final Output
$\boldsymbol{u}_{s}{ }^{*}=\operatorname{argmin} \mathrm{B}_{s}\left(\boldsymbol{u}_{s}\right)$

Step 3: Move and add Particles $\left\{\boldsymbol{u}_{s}\right\}$



Particles from $\boldsymbol{u}_{t}$

Step 5: Take $K$ best Particles $\left\{\boldsymbol{u}_{s}\right\}$


## Good since

$\psi_{s t}\left(\boldsymbol{u}_{s}, \boldsymbol{u}_{t}\right)=$ $\omega\left[\boldsymbol{u}_{s} \neq \boldsymbol{u}_{t}\right]$

## Toy Example - Shift 4.5



50 particles
Energy: 21959
Error: 0.4159
Random init


1 particle
Energy: 22593

Error: 0.3864
Random init


Ground Truth


Error: 0.8259

## Animation



50 particles
Energy: 21959
Error: 0.4159


1 particle
Energy: 22593
Error: 0.3864

## Number of Particles



## Results



## Extension: 6D Scene Flow

[M. Hornacek, A. Fitzgibbon, C. Rother, CVPR 2014]


Left RGBD image


Our result


Right RGBD image


Closest competitor
[Herbst, Ren, Fox, ICRA 2013]

## Extension: Reflections on Stereo

[R. Nair, A. Fitzgibbon, D. Kondermann, C. Rother, ICCV 2015]


depth

normal

depth

normal

reflection

reflectivity roughness

## Related Work

- Different Variant of [Kothapa et al.]. Run full BP and then re-sample particles (no augmentation) [Peng et al. ICML '11]
- Maintain diverse particle set [Pacheco, Zuffi, Black, Sudderth, ICML '14]
- None Particle-based Methods (rarely applied to Computer Vision due to runtime limitations):
- Nonparametric belief propagation [Sudderth et al. IEEE Intl. Conf. Acoustics, Speech, Signal 2010]
- Sparse forward-backward [Pal et al. ICASSP 2006]
- Kernel BP [Song et al. AlStats 2011]
- Stochastic belief propagate [N. Noorshams and M. J. Wainwright arxiv 2011]


## See separate slides in learning.pdf

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[^0]:    * Loopy Belief Propagation is not exact method, but explained here due to its relation to dynamic programming!

